

# Approximating the Permanent via Nonabelian Determinants

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# Determinant and Permanent

- Let  $A$  be a  $n \times n$  matrix with entries in  $\{0, 1\}$

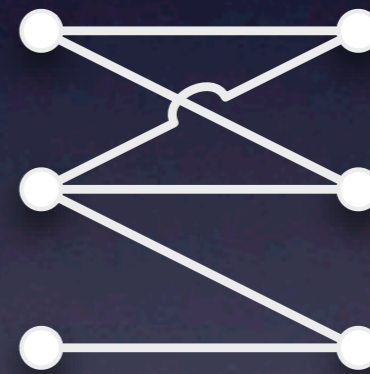
$$\det A = \sum_{\pi \in S_n} (-1)^\pi \prod_i A_{i, \pi i} \quad \text{perm } A = \sum_{\pi \in S_n} \prod_i A_{i, \pi i}$$

- $\det A$  : geometric meaning, basis-independent, homomorphic, easy
- $\text{perm } A$  : combinatorial, basis-dependent, hard

# A Combinatorial Picture

- Treat  $A$  as the adjacency matrix of a bipartite graph  $G$ :

$$\begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 0 & 1 \end{bmatrix}$$



- Then  $\text{perm } A = \#$  of *perfect matchings* in  $G$

# Complexity

- 0-1 PERMANENT is #P-complete [Valiant '79]
- However, it can be approximated using a rapidly-mixing Markov chain that samples random perfect matchings [Jerrum, Sinclair, Vigoda '04]
- Is there another approach, which is purely *algebraic*?

# The Godsil-Gutman Estimator

- Let  $A$  be a  $n \times n$  matrix with entries in  $\{0, 1\}$
- Let  $M_{ij} = \gamma_{ij} A_{ij}$  for uniformly random  $\gamma_{ij} \in \{\pm 1\}$
- Then

$$\mathbb{E} [(\det M)^2] = \text{perm } A$$

# What?!

- It's true and, better yet, simple:

$$\mathbb{E} [(\det M)^2] = \mathbb{E} \sum_{\pi, \sigma} (-1)^{\pi\sigma} \prod_i M_{i, \pi i} M_{i, \sigma i}$$

- only contributing terms appear when  $\pi = \sigma$  and all  $M_{i, \pi i}$  are nonzero:

$$= \sum_{\pi} (-1)^{\pi\pi} \left( \prod_i M_{i, \pi i} \right)^2 = \sum_{\pi} \prod_i A_{i, \pi i}$$

# But What's the Variance?

- For matrix  $A$ , define  $X = (\det M)^2$ . Then

$$\mathbb{E}[X] = \text{perm } A$$

- How many samples do we need? Chebyshev:

$$t \approx \frac{\mathbb{E}[X^2]}{\mathbb{E}[X]^2}$$

- Can we control this ratio?

# Sadly...

- [KKLLL '93]:  $\frac{\mathbb{E}[X^2]}{\mathbb{E}[X]^2} = 3^{n/2} \cdot \text{poly}(n)$
- But, they show that if  $\gamma_{ij}$  is uniform on the unit circle, or even just the cube roots of 1,

$$\frac{\mathbb{E}[X^2]}{\mathbb{E}[X]^2} = 2^{n/2} \cdot \text{poly}(n)$$

- Can we do better?



# Higher-dimensional algebras?

- Barvinok: what if the  $\gamma_{ij}$  are quaternions? Or higher-dimensional objects?
- [Chien, Rasmussen, Sinclair '03]  
In the *Clifford algebra* of dimension  $d$ :

$$\left(1 + O\left(\frac{1}{d}\right)\right)^{n/2}$$

- In particular, for the quaternions,  $\left(\frac{3}{2}\right)^{n/2}$

# How to define the determinant?

- In the nonabelian case, order matters
- The conventional determinant takes each product from top to bottom: no efficient algorithm is known!
- [Barvinok '00]: *symmetrize* each term:

$$\text{sdet } M = \frac{1}{n!} \sum_{\pi, \alpha \in S_n} (-1)^\pi \prod_i M_{\alpha i, \pi \alpha i}$$

- $O(n^d)$  algorithm in a  $d$ -dimensional algebra

# Algebraic Estimators

- Define  $M_{ij} = \rho_{ij} A_{ij}$ , where  $\rho_{ij}$  is drawn from some distribution on a nonabelian algebra  $\mathcal{A}$
- The Haar measure on unitary  $d \times d$  matrices
- The Gaussian measure on  $d \times d$  matrices (independent entries)
- $\det M$  takes values in  $\mathcal{A}$
- Define  $X = \|\det M\|^2$  or  $X = |\text{tr } \det M|^2$

# Our Results

- Two estimators, Haar or Gaussian measure:

$$X = |\operatorname{tr} \det M|^2 \quad X_s = |\operatorname{tr} \operatorname{sdet} M|^2$$

- We establish the following ratios:

$$\frac{\mathbb{E}[X^2]}{\mathbb{E}[X]^2} = \left(1 + O\left(\frac{1}{d}\right)\right)^n \quad \frac{\mathbb{E}[X_s^2]}{\mathbb{E}[X_s]^2} = \Omega\left(\frac{2^n}{n^d}\right)$$

- Ratios differ by  $O(d^4)$  for  $X = \|\det M\|^2$

# Expansion

- Let's expand the symmetric estimator:

$$X_s = \sum_{\kappa, \lambda \vdash A} (-1)^{\kappa\lambda} \mathbb{E}_{\alpha, \beta} \left( \text{tr} \prod_i \rho_{\alpha i, \kappa \alpha i} \right) \left( \text{tr} \prod_i \rho_{\beta i, \lambda \beta i}^* \right)$$

where  $\kappa \vdash A \Leftrightarrow \forall i : A_{i, \kappa i} = 1$

- Again, to contribute to the expectation,  $\kappa = \lambda$  so each  $\rho_{ij}$  appears twice or not at all

# Permuted Products

$$\begin{aligned}\mathbb{E}_{\{\rho_{ij}\}}[X_s] &= \sum_{\kappa \vdash A} \mathbb{E}_{\{\sigma_i\}} \mathbb{E}_{\alpha, \beta} \left( \text{tr} \prod_i \sigma_{\alpha i} \right) \left( \text{tr} \prod_i \sigma_{\beta i}^* \right) \\ &= a_d \cdot \text{perm } A\end{aligned}$$

- where

$$a_d = \mathbb{E}_{\{\sigma_i\}} \mathbb{E}_{\alpha, \beta} \left( \text{tr} \prod_i \sigma_{\alpha i} \right) \left( \text{tr} \prod_i \sigma_{\beta i}^* \right)$$

- Covariance between a product of the same matrices  $\sigma_i$  taken in two different orders

# The Cupcap Cometh

- What, for instance, is

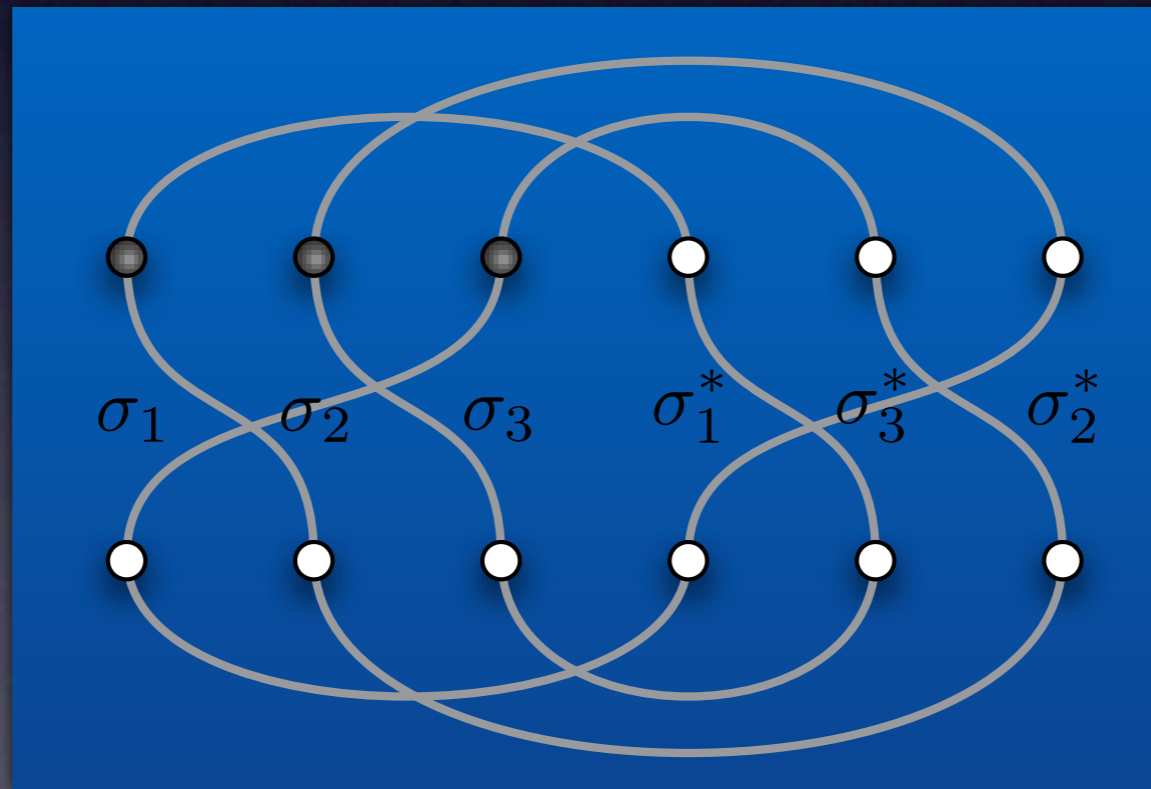
$$\mathbb{E}_{\sigma_1, \sigma_2, \sigma_3} (\text{tr } \sigma_1 \sigma_2 \sigma_3) (\text{tr } \sigma_1 \sigma_3 \sigma_2)^* ?$$

- Both Haar and Gauss:  $\mathbb{E}[\sigma_j^i (\sigma_\ell^k)^*] = \frac{1}{d} \delta^{ik} \delta_{j\ell}$
- Diagrammatically:

$$\mathbb{E}[\sigma \otimes \sigma^*] = \frac{1}{d} \begin{array}{c} \cup \\ \cap \end{array}$$

# Diagrams and Loops

- Form product by “weaving” matrices, and connect with cupcaps
- Tracing gives a factor of  $d$  for each loop



$$\mathbb{E}_{\sigma_1, \sigma_2, \sigma_3} (\text{tr } \sigma_1 \sigma_2 \sigma_3) (\text{tr } \sigma_1 \sigma_3 \sigma_2)^* = \frac{1}{d^2}$$



# A Generating Function

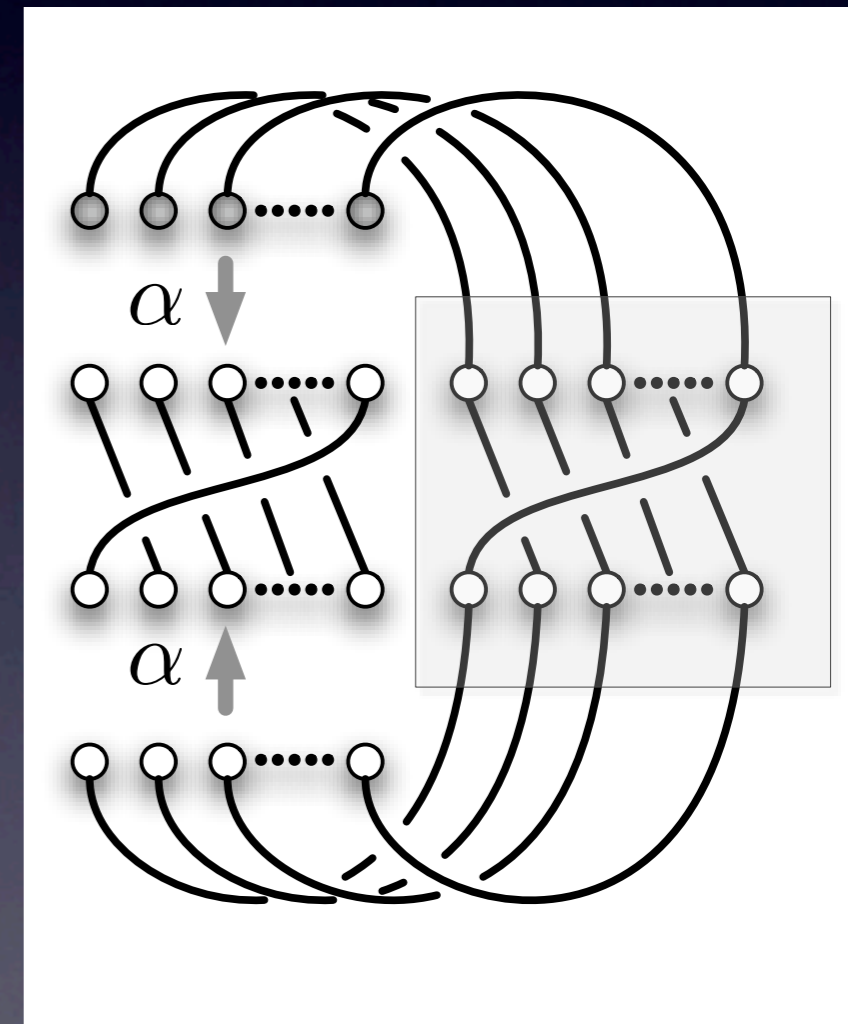
- Averaging over all permutations gives

$$a_d = \frac{1}{d^n} \mathbb{E}_\alpha d^{c([\alpha, r])}$$

- where  $c(\pi)$  is the number of cycles and  $r = (1\ 2\ \dots\ n)$  is a rotation

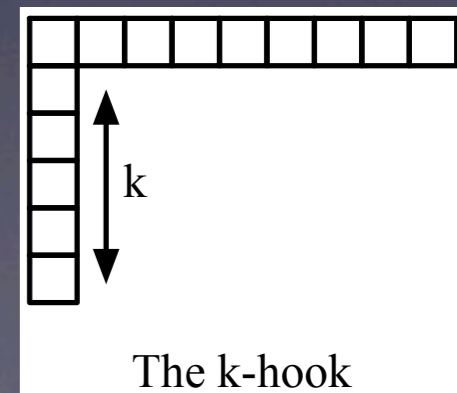
- Using Fourier analysis on  $S_n$ ,

$$a_d = \frac{1}{d^n} \left( \binom{n+d}{n+1} - \binom{d}{n+1} \right)$$



# Fourier Analysis

- First we write  $a_d = \frac{1}{d^n} \mathbb{E}_{r,r'} d^{c(rr')}$
- Then  $a_d$  is an inner product  $\langle d^{c(\cdot)}, P * P \rangle$   
 where  $P$  is the uniform distribution on  $n$ -cycles
- $d^{c(\cdot)}$  is the trace of a *combinatorial representation*:  
 action of  $S_n$  on strings of length  $n$  over  $\{1, \dots, d\}$
- Fourier coefficients are *Kostka numbers*
- $P$  is supported on “hooks”



# The Second Moment: Unsymmetrized

$$\mathbb{E}[X^2] = \sum_{C \vdash A} \sum_{(\kappa, \lambda, \mu, \nu) \vdash C} \mathbb{E}_{\{\rho_{ij}\}} \left( \text{tr} \prod_i \rho_{i, \kappa i} \right) \left( \text{tr} \prod_i \rho_{i, \lambda i} \right) \left( \text{tr} \prod_i \rho_{i, \mu i}^* \right) \left( \text{tr} \prod_i \rho_{i, \nu i}^* \right)$$

- The tuple  $(\kappa, \lambda, \mu, \nu)$  determines a *double cycle cover*: each  $\rho_{ij}$  appears 2 or 4 times

$$(\text{perm } A)^2 = \sum_{M_1 \vdash A} \sum_{M_2 \vdash A} = \sum_C 2^{|C|}$$

- $M_1 \oplus M_2 = C$ , a cycle cover

# Fourth-Order Operator

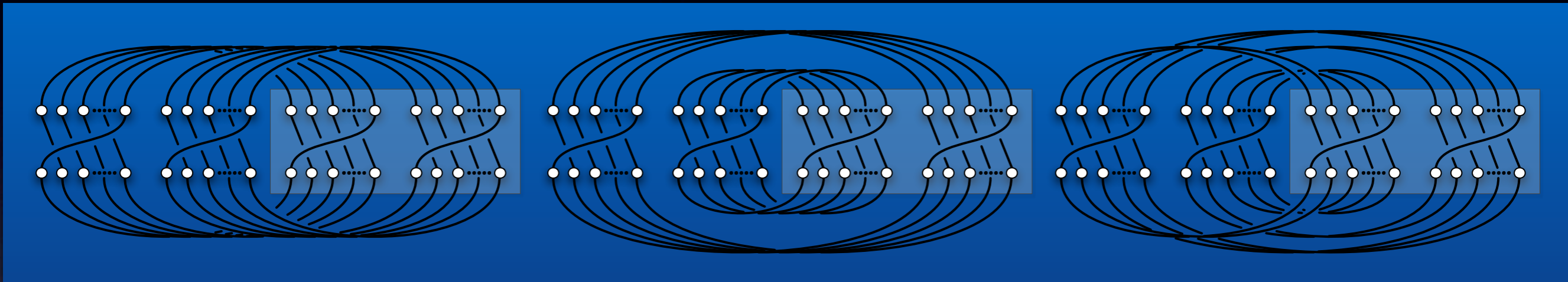
- In the Gaussian measure,

$$\mathbb{E}_\sigma [\sigma \otimes \sigma \otimes \sigma^* \otimes \sigma^*] = \frac{1}{d^2} \left( \begin{array}{c} \cup \\ \cup \\ \cup \\ \cup \end{array} + \begin{array}{c} \cup \\ \cup \\ \cup \\ \cup \end{array} \right)$$

- In the Haar measure,

$$\mathbb{E}_\sigma [\sigma \otimes \sigma \otimes \sigma^* \otimes \sigma^*] \preceq \left( 1 + O\left(\frac{1}{d}\right) \right) \frac{1}{d^2} \left( \begin{array}{c} \cup \\ \cup \\ \cup \\ \cup \end{array} + \begin{array}{c} \cup \\ \cup \\ \cup \\ \cup \end{array} \right)$$

# The Second Moment: Unsymmetrized



- Sum over all mixtures of (1,3),(2,4) and (1,4),(2,3) matchings
- Mixed matchings have fewer loops

$$2 \sum_{i=0,2,4,\dots}^n \binom{n}{i} d^{-i} = \left( \left(1 + \frac{1}{d}\right)^n + \left(1 - \frac{1}{d}\right)^n \right)$$

# Second Moment, Symmetrized

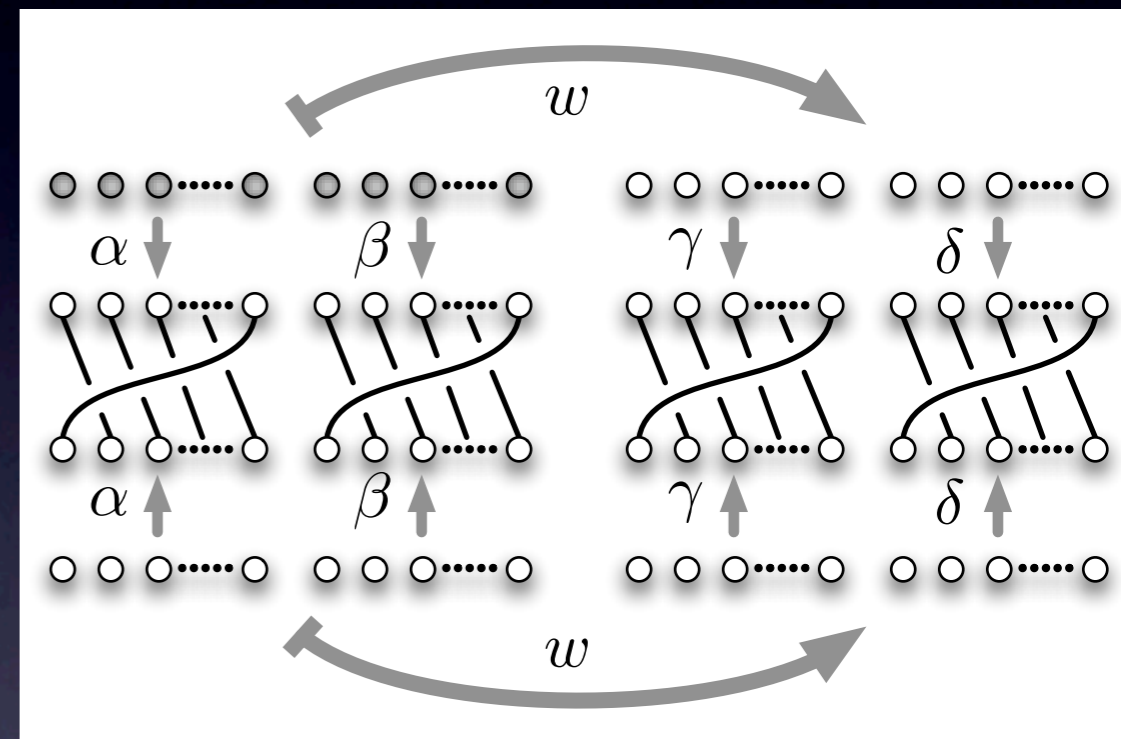
- Now the second moment can be bounded in terms of

$$\mathbb{E}_{\pi, \sigma \in S_{2n}} d^c(\pi^{-1}(r, r)\pi \sigma^{-1}(r, r)\sigma)$$

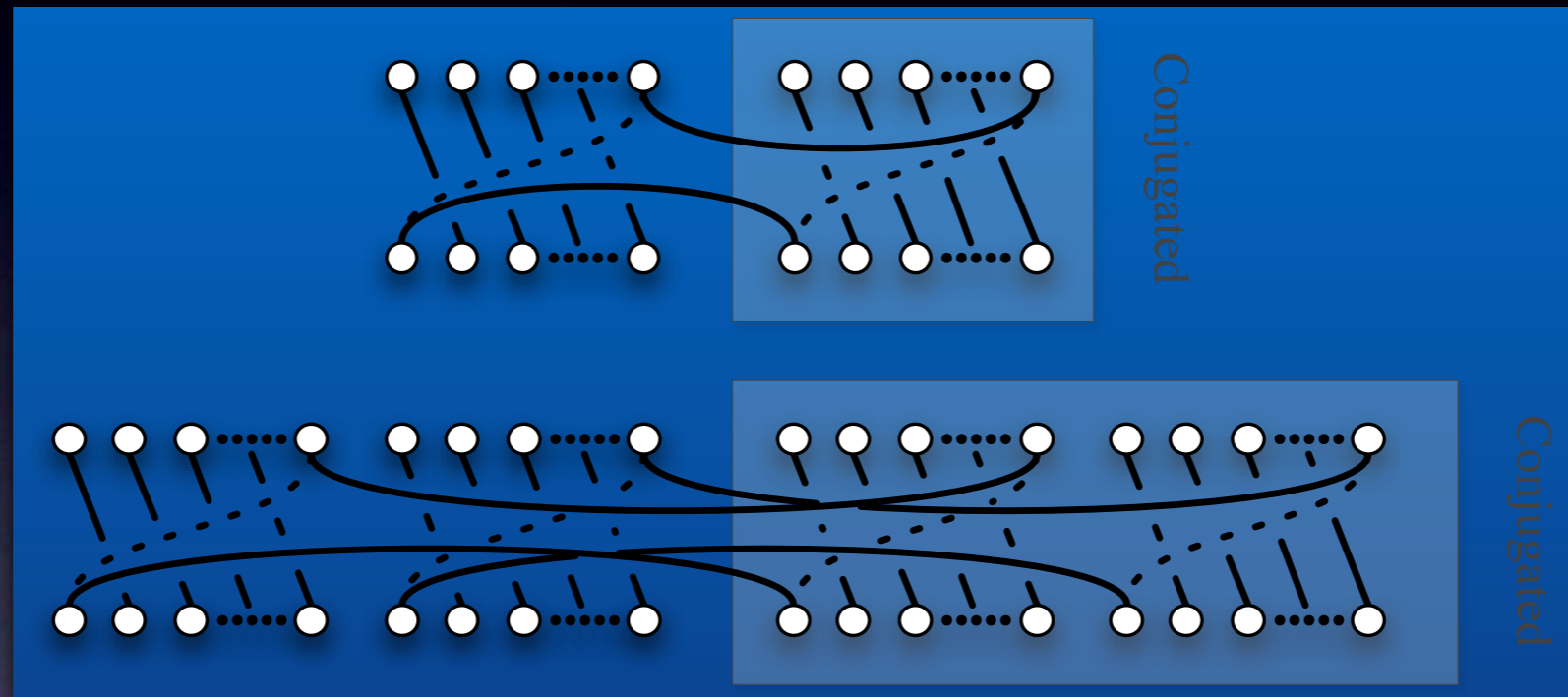
- Complicated distribution of pairs of  $n$ -cycles in  $S_{2n}$

- Bound in terms of uniform distribution

- Littlewood-Richardson rule: restrictions of irreps of  $S_{2n}$  to the Young subgroup  $S_n \times S_n$



# The Frobenius Estimators



- $\|M\|^2 = \text{tr } MM^\dagger$  instead of  $(\text{tr } M)(\text{tr } M^*)$
- Just rewire 2 or 4 edges

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## THE NATURE *of* COMPUTATION



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