

Pick up handout at front of hall

10.45 start

Induced E.M.F.
 $e = Hlv \times 10^{-8}$
 $e \cdot l = H \times v \times l \times 10^{-8}$
 where $l =$ length of one coil side

$\therefore e = 2Hlv \sin \theta \times 10^{-8}$
 $= 2Hl \times \pi d n \sin \theta \times 10^{-8}$
 where $n = \frac{\text{revs.}}{\text{sec}}$
 $d = \text{diam. of coil}$

If the coil rotates the current is collected by a slip-ring

$e_{\text{max}} = E = 2Hl \times \pi d n$
 or $E = 2\pi \phi n S$
 for a coil of n turns

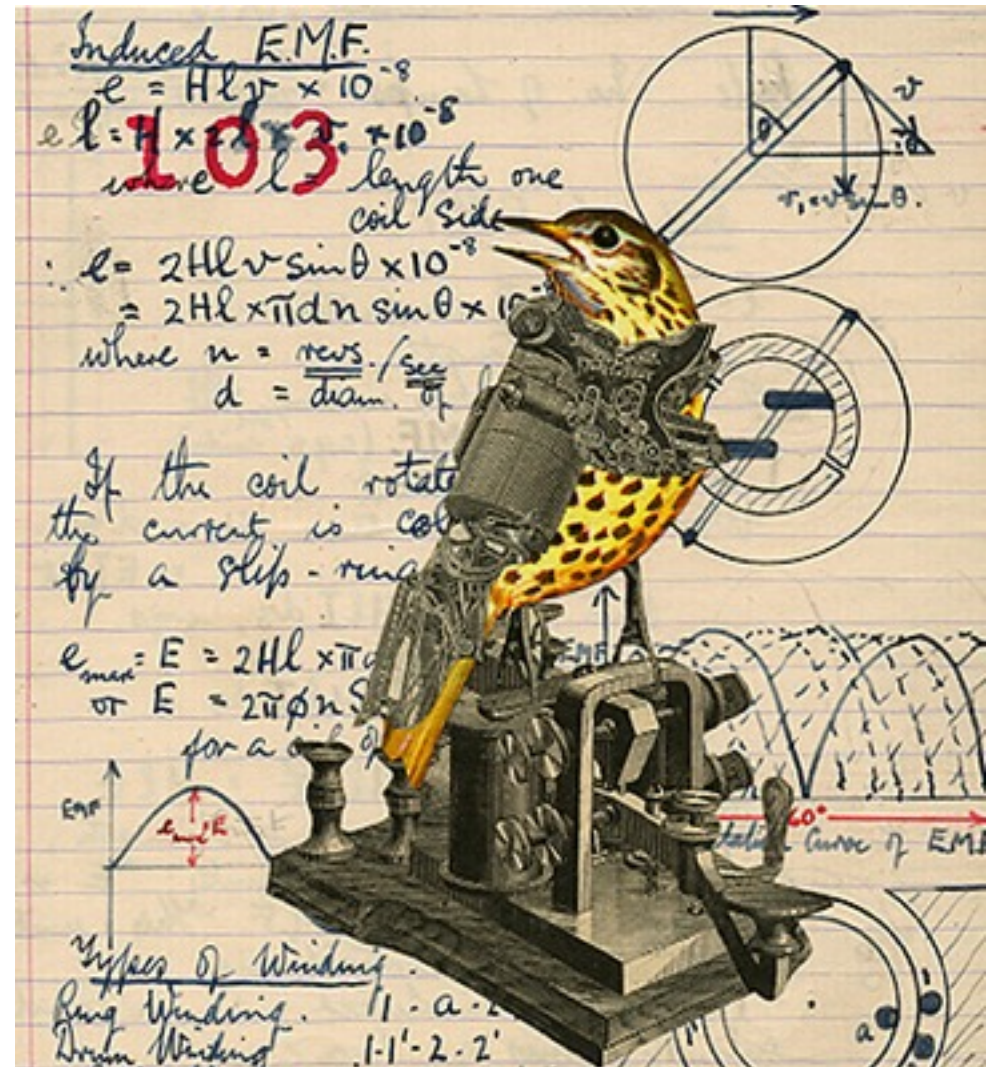
Graph of EMF vs. time showing a sine wave with amplitude e_{max} . A phase angle of 60° is indicated.

Types of Winding
 Lap Winding. 1-1'-2-2'
 Drum Winding. 1-1'-2-2'

Emergence

Automata Decomposition, Innovation, and Dissipation

Lecture 4
10:45 AM
28 June 2011



Topics

Hierarchies of Computation

The Chomsky Hierarchy

Into the Labyrinth

Finite State Automata

Non-deterministic vs. Probabilistic

Groups

Review of Group and Semigroup Theory

Jordan-Hölder Decomposition

Semigroups and Automata

The Amazing Krohn-Rhodes Theorem

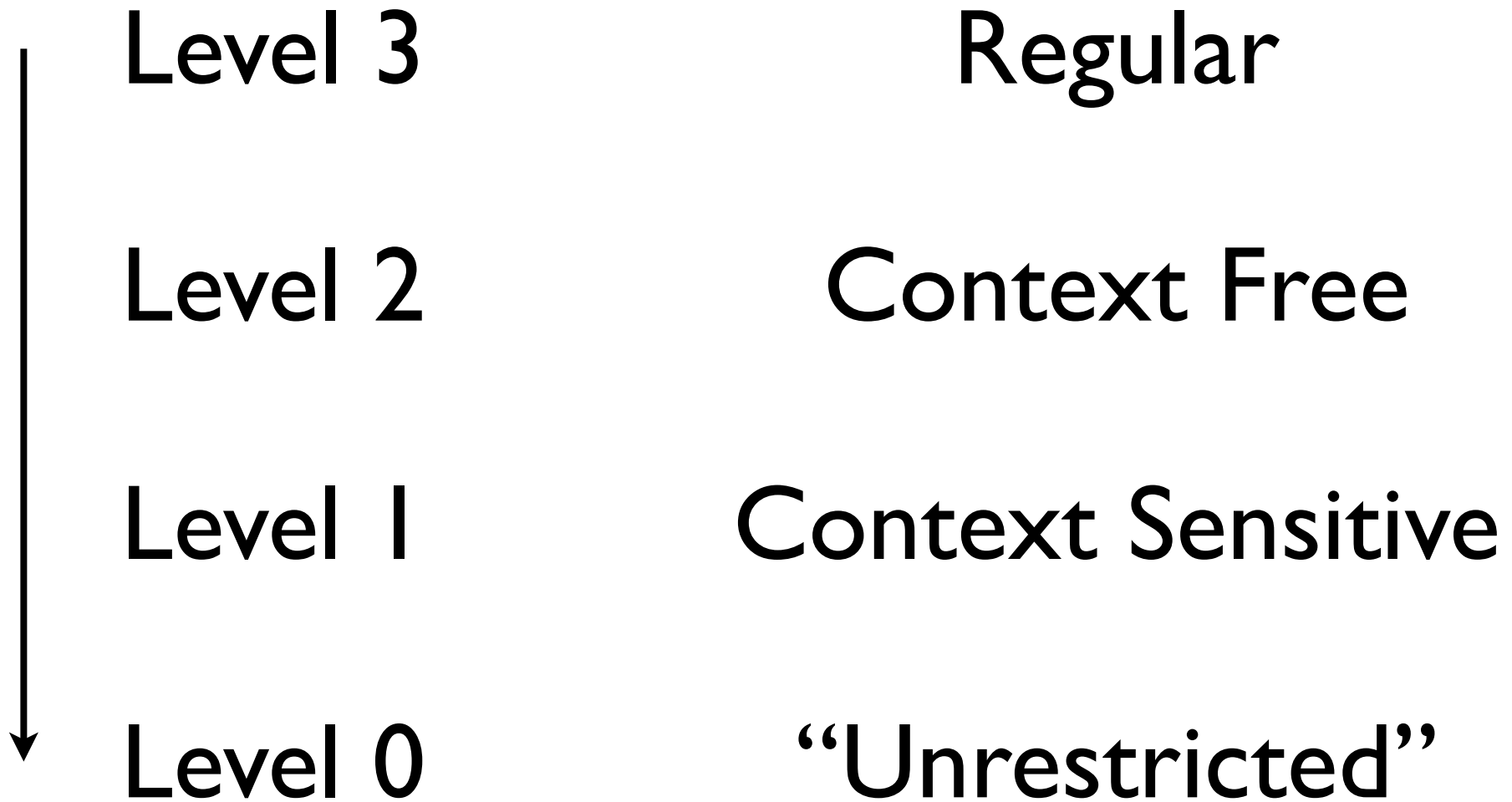
How to “Smooth” a Computer

Universal vs. Finite State Computation

Introducing Stochasticity

Prime Decomposition of the Bengalese Finch

The Chomsky Hierarchy



Jim's Bestiary

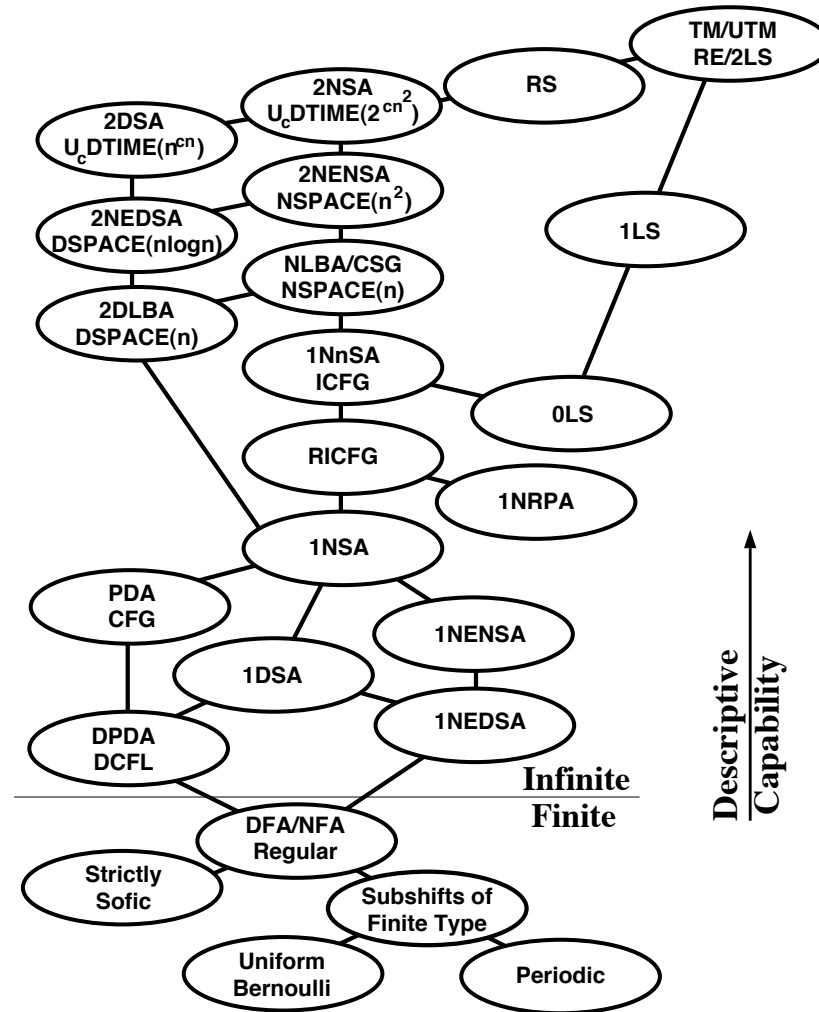
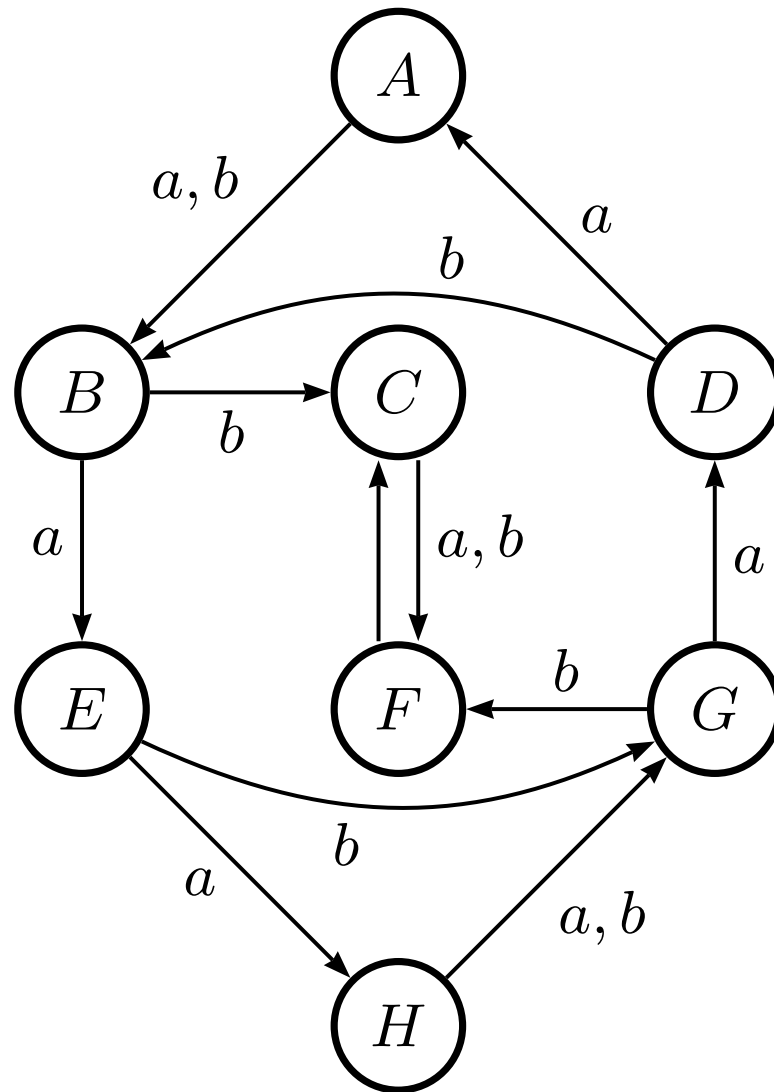


Figure 2 The discrete computation hierarchy. *Adjective legend:* 1 = one way input tape, 2 = two way input tape, D = deterministic, N = nondeterministic, I = indexed, RI = restricted I, n = nested, NE = nonerasing, CF = context free, CS = context sensitive, R = recursive, RE = R enumerable, and U = universal. *Object legend:* G = grammar, A = automata, FA = finite A, PDA = pushdown A, SA = stack A, LBA = linear bounded A, RPA = Reading PDA, TM = Turing machine, LS = Lindenmayer system, 0L = CF LS, 1L = CS LS, and RS = R set. (After [31,35–39].)

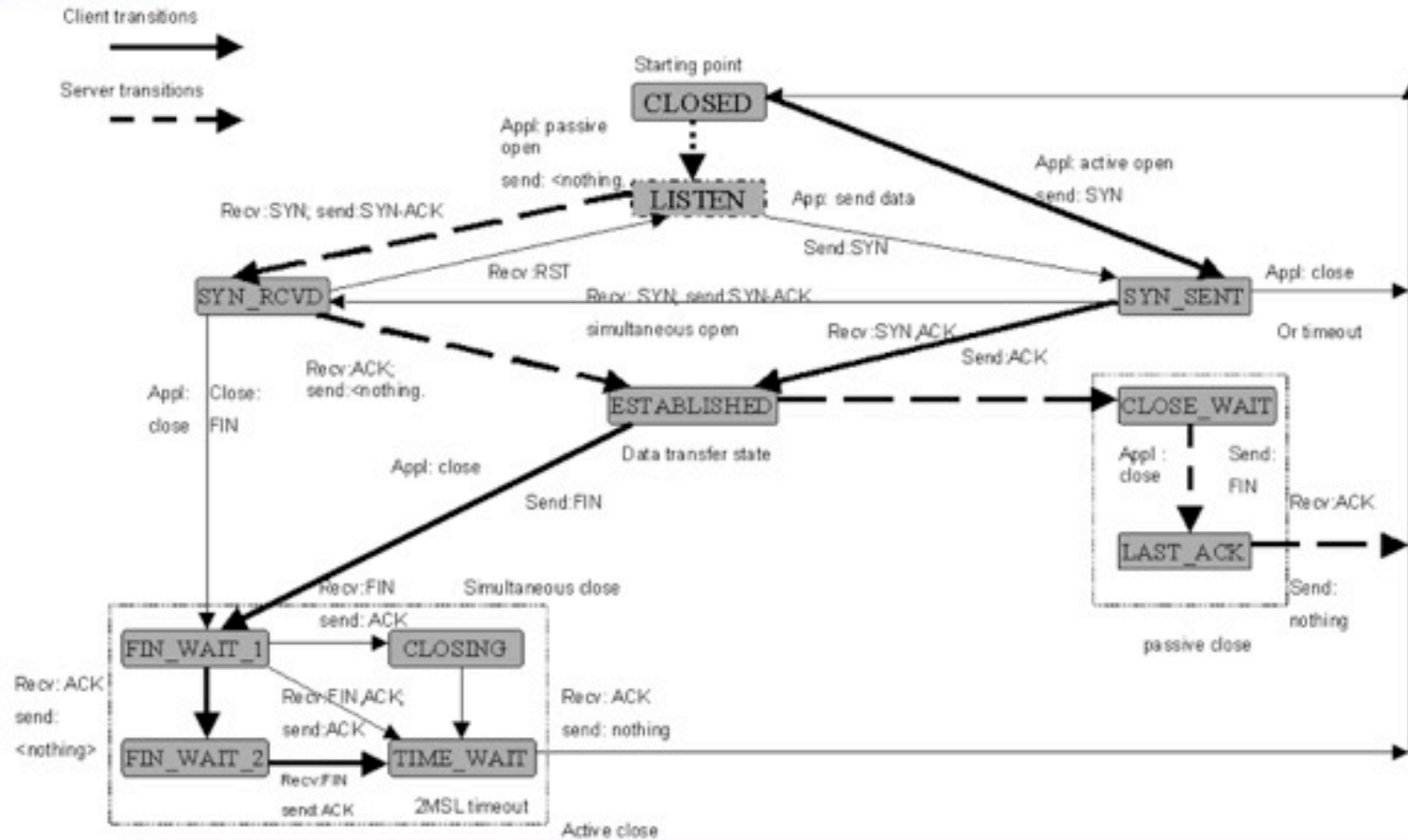
Finite State Automata



Finite State Automata



TCP State Machine (TCP/IP Illustrated vol. 1) W. Richard Stevens



Finite State Automata

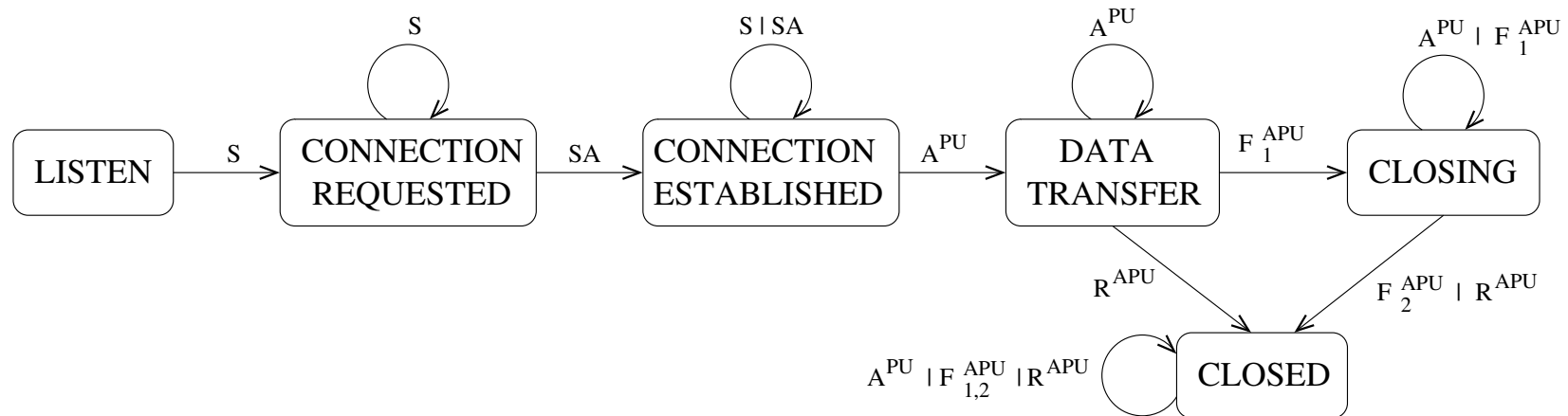
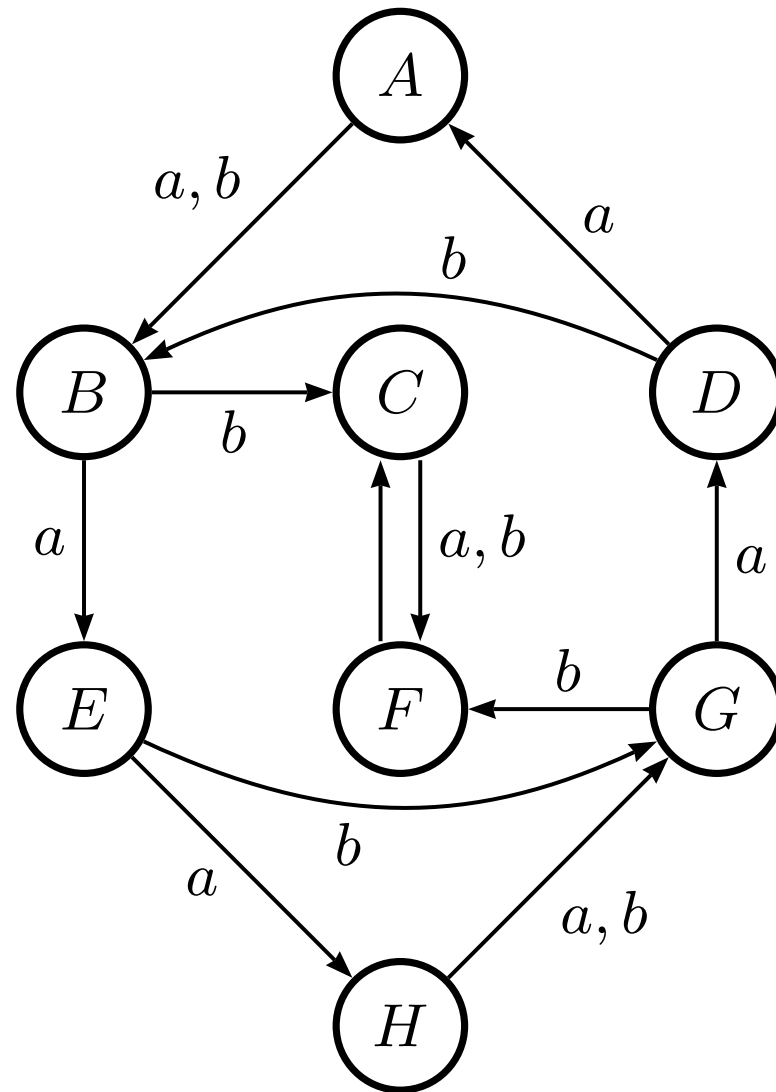


Figure 2: Graphical representation of the TCP FSM. The events are the flag combinations of Table 4 and the states are as defined in Table 3. For clarity, the failure state is not shown.

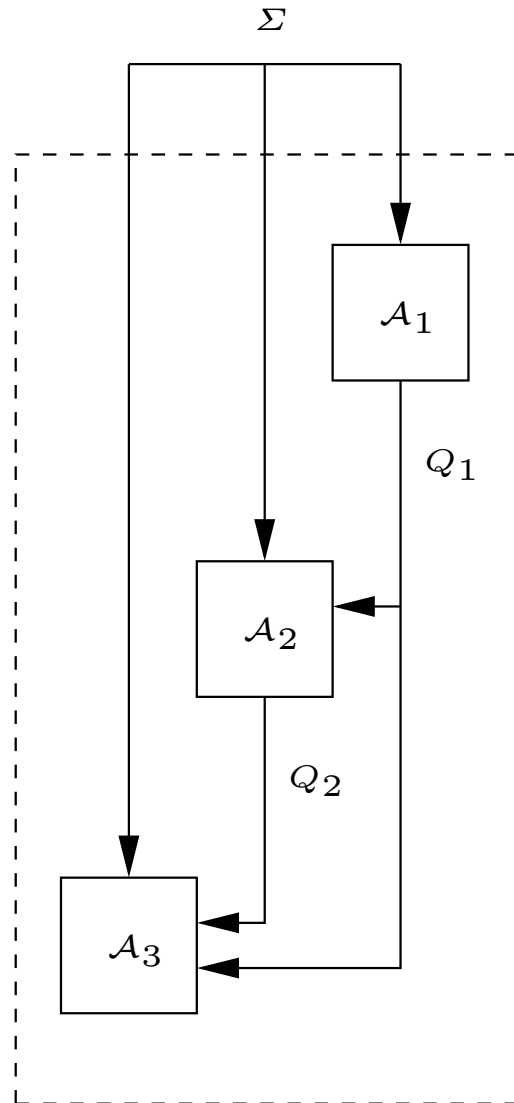
Abstract

Finite state machines can be used to detect anomalous behaviour in TCP traffic by describing the progression of a connection through states as a result of events based on header flags. The method was applied to real traffic to understand its realistic use and it was found that for the time period analysed here, on the order of 37% of TCP connections do not follow the TCP protocol specifications. The majority of these are a result of malicious activity, and approximately 4% are due to benign anomalies such as unresponsive hosts and misconfigurations. The method may be applied as a network security measure, as a network management tool or as a research tool for the study of TCP behaviour on the Internet.

Finite State Automata



Hierarchical Decomposition



Factoring the Z4 Machine

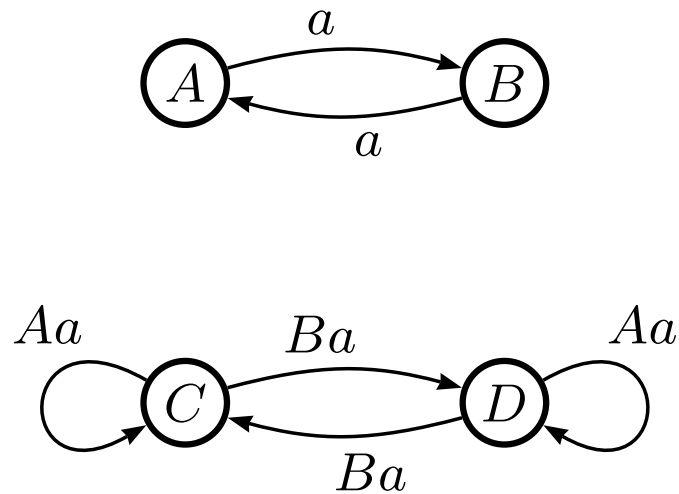
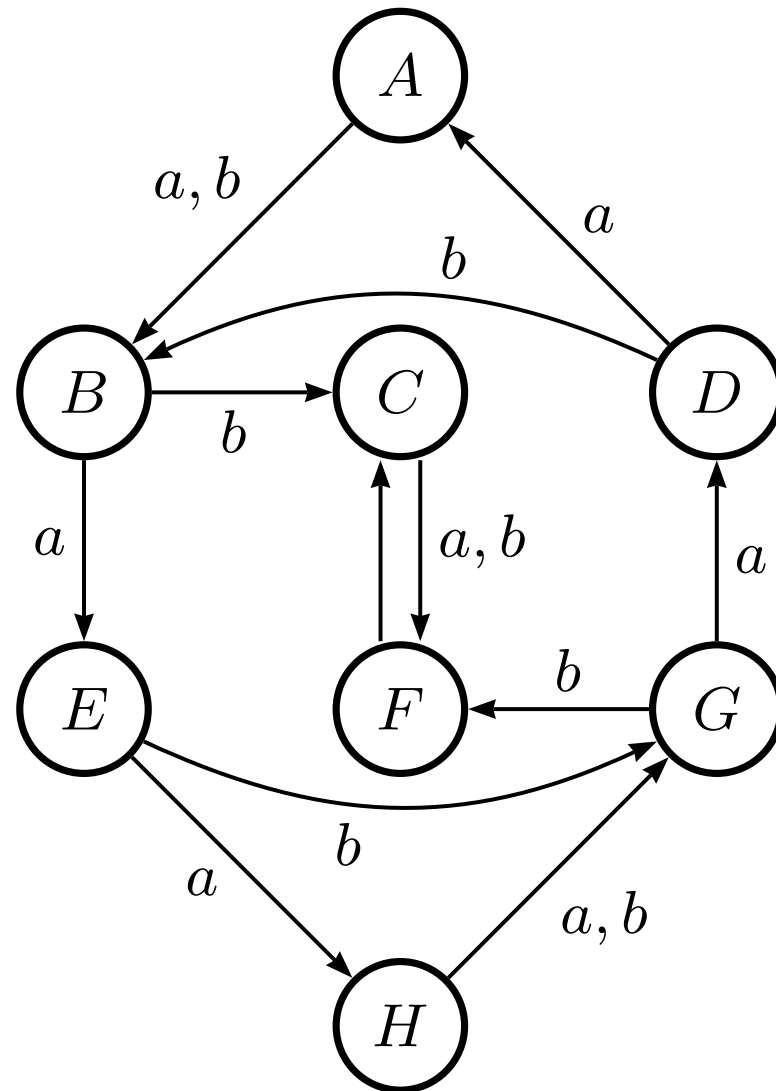
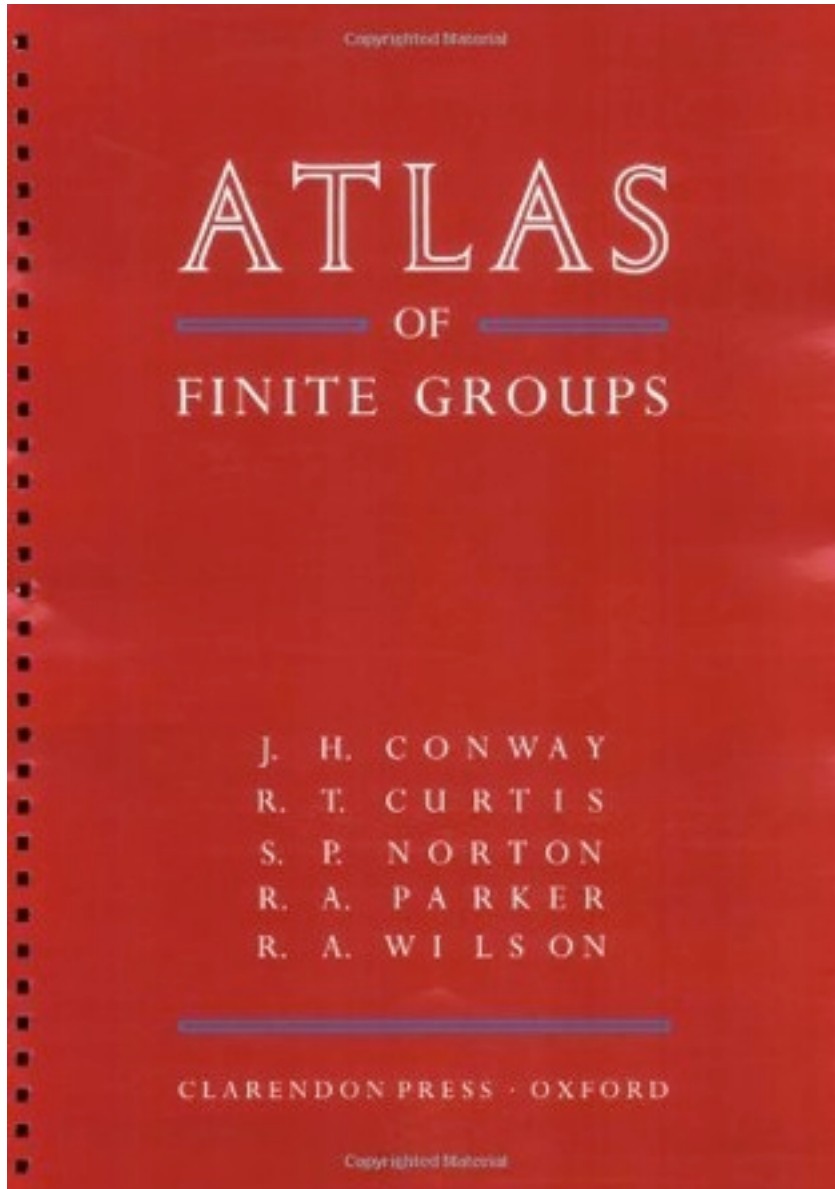


FIG. 3. A cascade of two Z_2 counters; the lower-level machine has two possible transitions, depending on the state of the higher-level machine. The cascade now can count modulo four, with the high-level machine counting even versus odd, and the lower-level machine tracking the second bit.

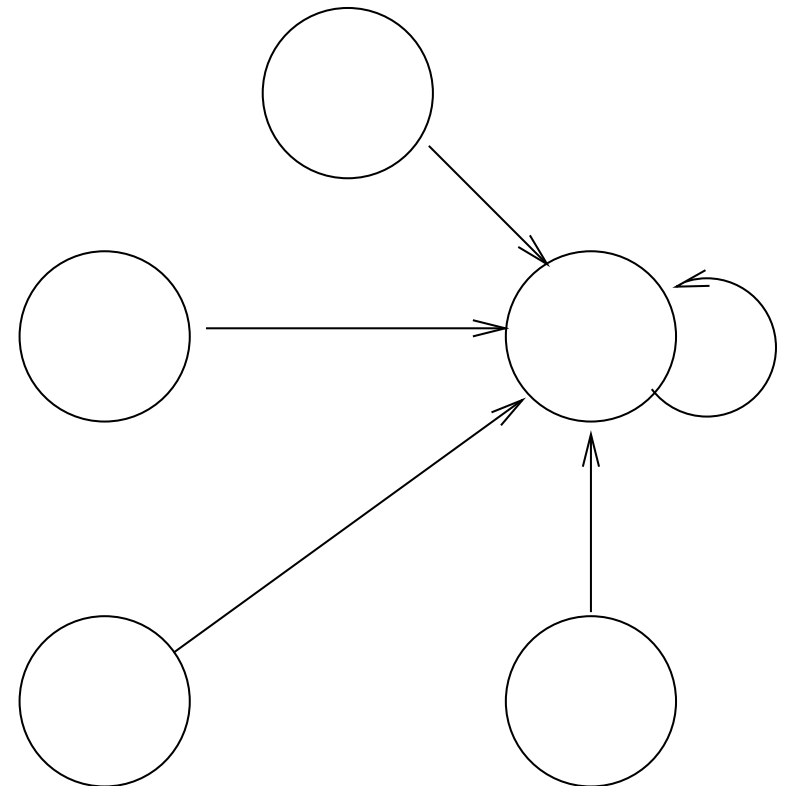
Finite State Automata



Krohn-Rhodes : *all* FSA are



+ “Pure” Resets



An Example of the Krohn-Rhodes Holonomy Decomposition

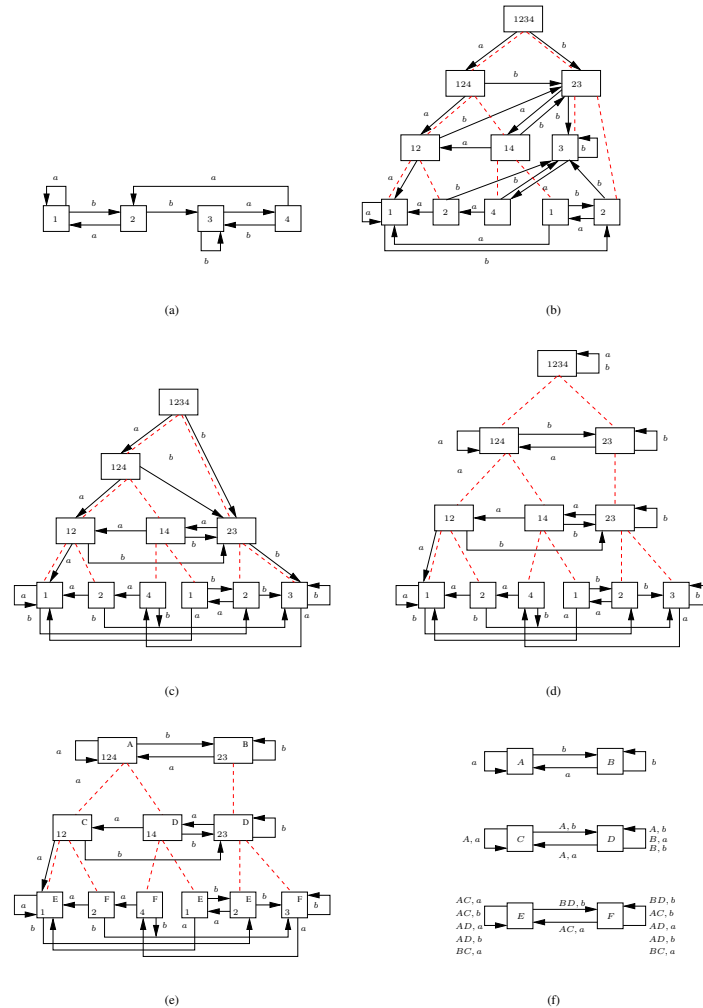


Fig. 6. The decomposition process: (a) An automaton; (b) its TSA (parenthood indicated by dashed lines); (c) the TSA rearranged according to height; (d) the holonomy tree obtained after completion and redirection; (e) state encoding; (f) the decomposition.

Another Example

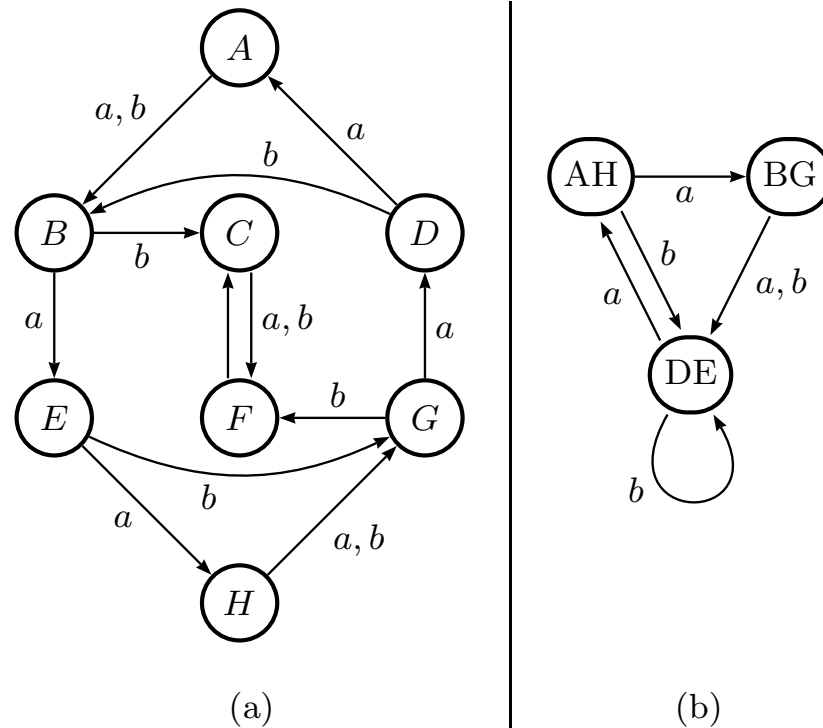
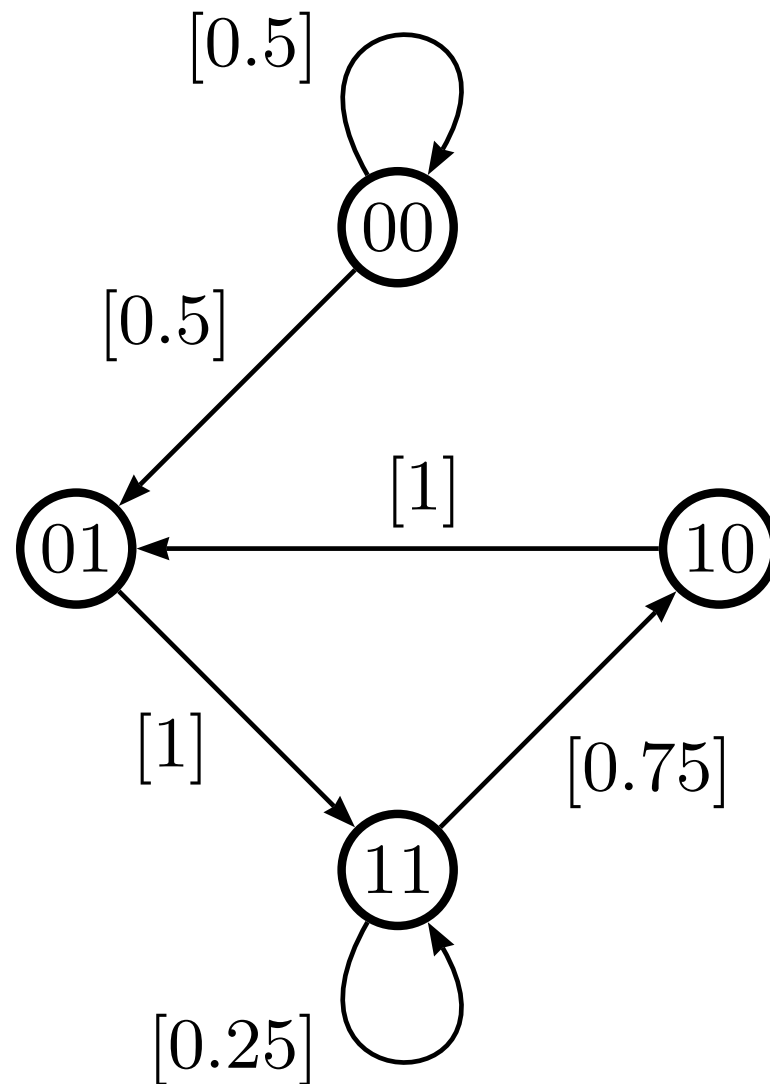


FIG. 4. Effective theories and “smoothed” computations. (a) An eight-state automaton with two input letters, built from XOR gates on the left-shift register. (b) The top level of the decomposed machine. Despite the complicated internal structure of (a), at the coarse-grained level, the process is seen count modulo three (on receiving a signals), and reset (on receiving b signals.) Movement within the three superstates of (b) is dictated by Z_2 counters and resets at lower levels in the decomposition.

Stochasticity — Random Transitions



Stochasticity — Random Transitions

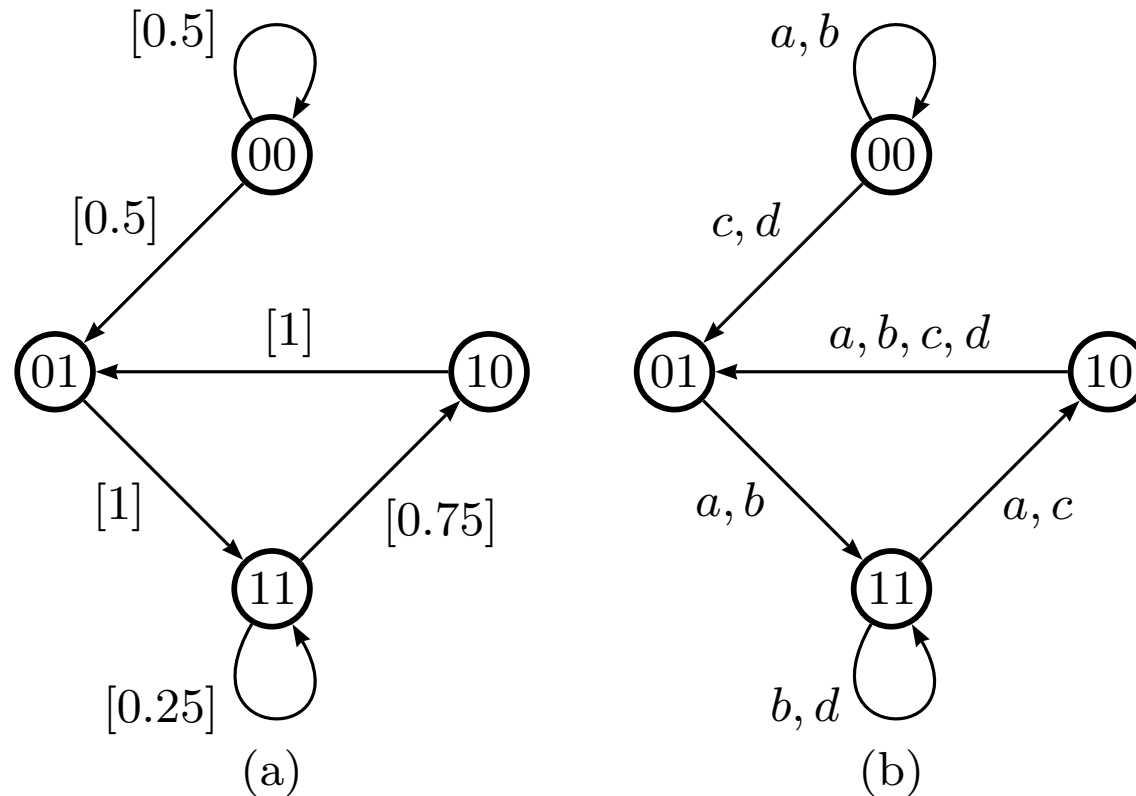


FIG. 5. Decomposition of a (single input letter) probabilistic automaton into a 4-letter deterministic machine; $\pi(a) = \pi(c) = 3/8$; $\pi(b) = \pi(d) = 1/8$.

Ashes to Ashes

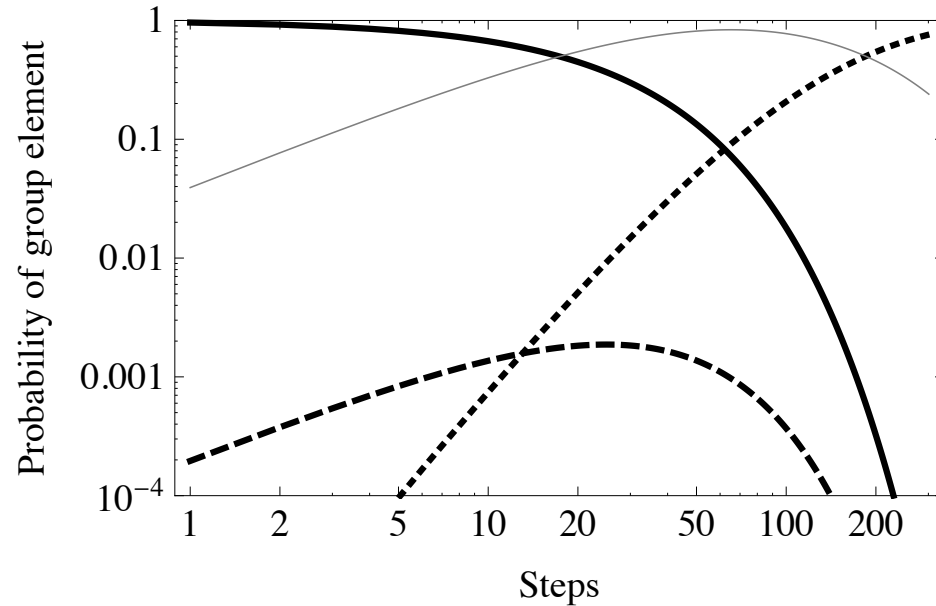


FIG. 6. The rise, and fall, of group structures in noisy Boolean systems. For the XOR function of Fig. 2b, with a 1% bit-flip rate, we show the probability that the function implements one of the group elements of the noise free case (Z_3 , solid line), one of the four reset automata (dotted line), or any group element not in Z_3 (the set $S_4 - Z_3$, dashed line.) The light gray line shows the remainder. As time passes and errors accumulate, the chances of a transformation being a member of Z_3 decline, while the reset automata come to dominate the long term input-output maps. The non-trivial relationship between noise and complexity can be seen when, for a brief period, representatives of more complex groups appear and even dominate the pure resets.

The Prime Decomposition of the Bengalese Finch

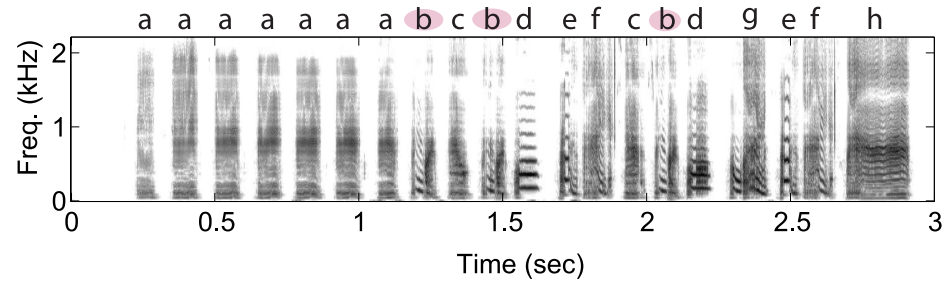


[Finch movie]

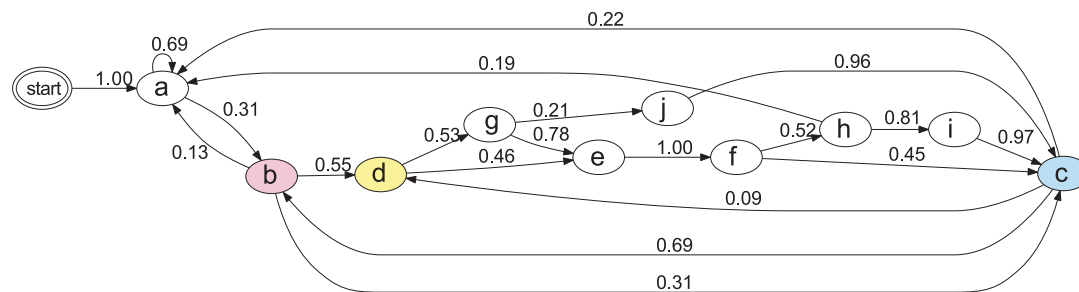
Complex sequencing rules of birdsong can be explained by simple hidden Markov processes

Kentaro Katahira, Kenta Suzuki, Kazuo Okanoya, Masato Okada
(<http://arxiv.org/abs/1011.2575>)

A

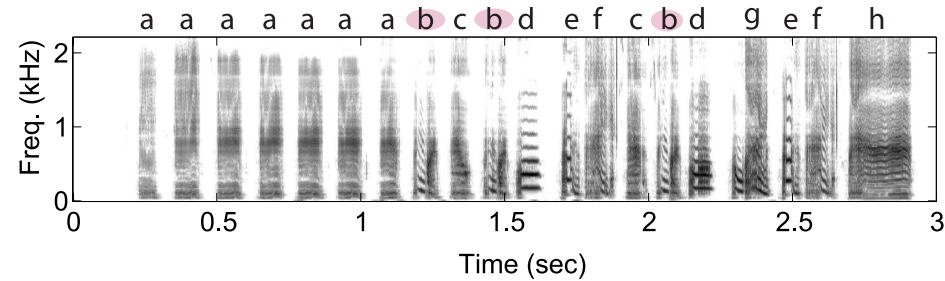


B

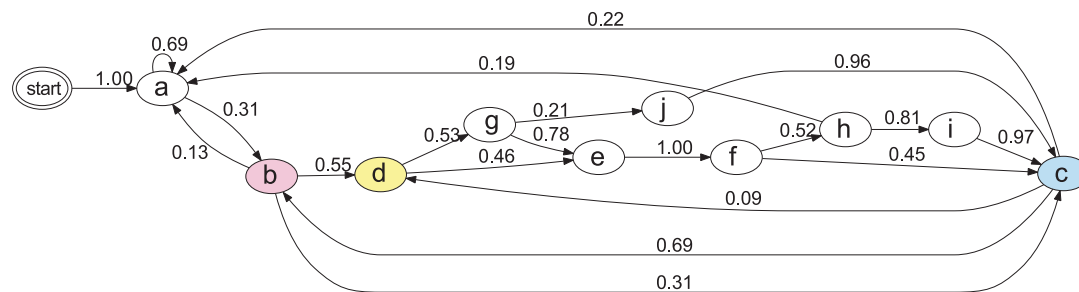


The Prime Decomposition of the Bengalese Finch

A



B



10 states $\rightarrow 10^{10}$ elements in the full transformation semigroup?